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THE FAMILY OF PDOL GROWTH-SETS  
IS PROPERLY INCLUDED IN THE FAMILY OF  
DOL GROWTH-SETS

BY

JUHANI KARHUMÄKI

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## 1. Introduction

L-systems have been introduced for biological purposes (see [3]). However, these have been studied intensively during the last few years from the formal language point of view (see [1] and [6]). A particularly interesting aspect within L-systems is the theory of growth functions. These have been studied for instance in [7] and [8].

The purpose of this note is to show that the family of PDOL growth-sets is properly included in the family of DOL growth-sets. As a corollary of this result we also solve a problem introduced in [5], namely that  $\mathcal{L}_{\text{CPDOL}} \subsetneq \mathcal{L}_{\text{CDOL}}$ .

## 2. Notations

We assume that the reader is familiar with the standard formal language notations. For the definitions of DOL-systems, -languages, and -sequences we refer to [1]. We say that a DOL-system  $G$  is  $\lambda$ -free, or a PDOL-system, iff there are no  $\lambda$ -productions in  $G$ .

If  $G$  is a DOL-system, then  $L(G)$  (resp.  $E(G)$ ) means the language (resp. the sequence) generated by  $G$ . The *growth-set* generated by  $G$  is

$$|L(G)| = \{|P| \mid P \in L(G)\},$$

where  $|P|$  means the length of the word  $P$ . Let

$$E(G) = \omega_0, \omega_1, \dots$$

Then the *growth-sequence* generated by  $G$  is

$$|E(G)| = |\omega_0|, |\omega_1|, \dots$$

We say that a homomorphism  $h : V_1^* \rightarrow V_2^*$  is a *coding* iff it maps each a letter to another letter. So all codings are length preserving.

A language  $L$  is called a CDOL-language (resp. a CPDOL-language) iff there exists a DOL-system  $G$  (resp. a PDOL-system  $G$ ) and a coding  $h$  such that

$$L = h(L(G)).$$

The family of CDOL-languages (resp. CPDOL-languages) is denoted by  $\mathcal{L}_{\text{CDOL}}$  (resp.  $\mathcal{L}_{\text{CPDOL}}$ ).

### 3. Lemmas

We need the following three lemmas.

**Lemma 1.** *Let*

$$(1) \quad |E(G)| = |\omega_0|, |\omega_1|, \dots$$

*be a DOL growth-sequence. Then (1) is ultimately periodic modulo 2.*

*Proof.* It is well known (see [7]), that, for  $n \geq n_0$ , (1) satisfies a recursion formula with integer coefficients, say

$$|\omega_n| = \sum_{i=1}^k \alpha_i |\omega_{n-i}|, \quad \text{for } n \geq n_0.$$

Let  $g$  be the canonical homomorphism of  $Z$  onto  $Z_2$ . By applying  $g$  to the above equation we obtain in the finite set  $Z_2$  the recursion formula

$$g(|\omega_n|) = \sum_{i=1}^k g(\alpha_i) g(|\omega_{n-i}|), \quad \text{for } n \geq n_0.$$

Thus, the sequence determined by this recursion formula must be ultimately periodic in  $Z_2$ . So we have proved Lemma 1.

Let  $H$  be the following DOL-system. The axiom is  $a$ , the alphabet is

$$V = \{a, b, c, a_1, b_1, c_1, a_2, b_2, c_2\},$$

and the productions are as follows:

$$\begin{aligned} a &\rightarrow a_1 a_2, & a_1 &\rightarrow abc^2, & a_2 &\rightarrow \lambda, \\ b &\rightarrow b_1 b_2, & b_1 &\rightarrow bc^2, & b_2 &\rightarrow \lambda, \\ c &\rightarrow c_1 c_2, & c_1 &\rightarrow c, & c_2 &\rightarrow \lambda. \end{aligned}$$

Let

$$L_1 = \{P \in V^* | a \Rightarrow^{2^n} P, \text{ for some } n \geq 0\}$$

and

$$L_2 = \{P \in V^* | a \Rightarrow^{2^{n+1}} P, \text{ for some } n \geq 0\}.$$

Then the corresponding sequences are

$$E_1 = a, abc^2, abc^2bc^4, abc^2bc^4bc^6, \dots$$

and

$$E_2 = a_1 a_2, a_1 a_2 b_1 b_2 (c_1 c_2)^2, a_1 a_2 b_1 b_2 (c_1 c_2)^2 b_1 b_2 (c_1 c_2)^4, \dots$$

So the language generated by  $H$  is

$$L(H) = L_1 \cup L_2 = \{a, abc^2bc^4 \dots bc^{2n} | n \geq 1\} \cup \{h(P) | P \in L_1\},$$

where  $h$  is the homomorphism of  $\{a, b, c\}$  into  $V$  defined by  $h(y) = y_1 y_2$ . Because

$$1 + n + 2 + 4 + \dots + 2n = (n + 1)^2$$

the growth-set determined by  $H$  is

$$|L(H)| = \{n^2, 2n^2 | n \geq 1\}.$$

We now put the elements of  $|L(H)|$  in increasing order and let  $X$  be this sequence. Denote

$$(2) \quad X = x_1, x_2, \dots$$

**Lemma 2.**  $X$  is not a DOL growth-sequence.

*Proof.* Assume the contrary: that a DOL-system  $H_1$  generates the sequence  $X$ . Then, by Lemma 1, (2) is ultimately periodic modulo 2. So there exist natural numbers  $r$  and  $s$  such that

$$x_r \text{ is an odd square}$$

and for each  $i$  and  $j$ ,  $i \geq 0$ ,  $0 \leq j \leq s - 1$ ,

$$x_{r+j+is} \equiv x_{r+j+(i+1)s} \pmod{2}.$$

Let  $x_r = k^2$ . Note that in (2) all odd integers are squares. So if  $m$  is the number of odd integers in the period, then for each  $i \geq 0$

$$x_{r+is} = (k + i2m)^2.$$

But this implies that in each period there must also be a fixed number of integers of the form  $2n^2$ . This however leads to a contradiction, as we shall now show.

For all  $i \geq 0$ , consider the natural numbers  $z_i$  satisfying the condition

$$(k + i2m)^2 < 2z_i^2 < (k + (i + 1)2m)^2,$$

or equivalently,

$$(3) \quad \frac{k}{\sqrt{2}} + i\sqrt{2}m < z_i < \frac{k}{\sqrt{2}} + i\sqrt{2}m + \sqrt{2}m.$$

By what we have shown, the number of such  $z_i$ 's is the same for all  $i \geq 0$ .

Trivially this number is either  $[\sqrt{2} m]$  or  $[\sqrt{2} m] + 1$ . Let  $\delta_0$  and  $\delta_1$  be positive real numbers defined by

$$\delta_0 = \sqrt{2} m - [\sqrt{2} m]$$

and

$$\delta_1 = [\sqrt{2} m] + 1 - \sqrt{2} m.$$

First, assume that for all  $i \geq 0$  the number of  $z'_i$ 's is  $[\sqrt{2} m]$ . Choose  $i_0$  such that  $i_0 \delta_0 > 1$ . Then the length of the interval

$$\left( \frac{k}{\sqrt{2}}, \frac{k}{\sqrt{2}} + i_0 \sqrt{2} m \right)$$

is

$$i_0 \sqrt{2} m = i_0 [\sqrt{2} m] + i_0 \delta_0 > i_0 [\sqrt{2} m] + 1.$$

So the number of  $z'_i$ 's in this interval is at least  $i_0 [\sqrt{2} m] + 1$ . On the other hand, by our assumption, their number is  $i_0 [\sqrt{2} m]$ .

Secondly, assume that the number of  $z'_i$ 's is  $i_0 [\sqrt{2} m] + 1$ , for all  $i \geq 0$ . Now choose  $i_1$  such that  $i_1 \delta_1 > 2$ . Thus, the length of the interval

$$\left( \frac{k}{\sqrt{2}}, \frac{k}{\sqrt{2}} + i_1 \sqrt{2} m \right)$$

is

$$i_1 \sqrt{2} m = i_1 [\sqrt{2} m] + i_1 - i_1 \delta_1 < i_1 [\sqrt{2} m] + i_1 - 2.$$

Thus, in this interval there are at most  $i_1 [\sqrt{2} m] + i_1 - 1$  numbers  $z_i$ . But by our assumption, in this interval there must be  $i_1 [\sqrt{2} m] + i_1$  numbers of this kind.

Because both the cases lead to a contradiction, we have proved Lemma 2.

We also need the following lemma, which is Lemma 5.4. of [4].

**Lemma 3.** *Let  $G$  be a PDOL-system generating an infinite language. Then there exists a PDOL-system  $G_1$  such that*

- (i)  $|L(G)| = |L(G_1)|$ ,
- (ii) *the sequence  $|E(G_1)|$  is strictly increasing.*

#### 4. Results

Now we are ready to establish our results.

**Theorem 1.** *The family of PDOL growth-sets is a proper subset of the family of DOL growth-sets.*

*Proof.* The inclusion is trivial. It is proper because  $|L(H)|$  is a DOL growth-set, but is not, by Lemmas 2 and 3, a PDOL growth-set.

As an immediate corollary of Theorem 1 we can solve a problem proposed in [5].

**Theorem 2.** *The family  $\mathcal{L}_{\text{CPDOL}}$  is properly included in the family  $\mathcal{L}_{\text{CDOL}}$ .*

Examples of languages which lie in the difference  $\mathcal{L}_{\text{CDOL}} \setminus \mathcal{L}_{\text{CPDOL}}$  are the languages  $L(H)$  and  $L = \{a^{n^2}, a^{2n^2} | n \geq 1\}$ .

We can generalize Theorem 1 to cover all *growing* DOL-systems, i.e., systems with an increasing growth-sequence. Of course any PDOL-system is a growing DOL-system.

**Theorem 3.** *The family of growth-sets generated by growing DOL-systems is properly included in the family of DOL growth-sets.*

*Proof.* It suffices to show that Lemma 3 can be generalized for growing DOL-systems. Assume that  $G$  is a growing DOL-system with  $E(G) = \omega_0, \omega_1, \dots$ . In the following we use the notations of [4].

Let  $M$  be the growth-matrix of  $G$ ,  $\pi$  the Parikh-vector associated with the axiom of  $G$ , and  $\eta$  the column vector with all elements equal to 1. Then the sequence

$$(4) \quad d_n = \pi(M^n - M^{n-1})\eta = \pi(M - I)M^{n-1}\eta, \quad n = 1, 2, \dots$$

tells us how much the length of the word grows during the  $n$ th step of the derivation. By (4), the  $d_n$ 's satisfy a recursion formula with integer coefficients. Thus, by the Theorem proved in [2], zeros occur in (4) ultimately periodically.

For all  $i \geq 0$ , let  $\text{Min}(\omega_i)$  denote the set of symbols occurring in  $\omega_i$ . It is well known that the sequence

$$(5) \quad \text{Min}(\omega_0), \text{Min}(\omega_1), \dots$$

is ultimately periodic.

Consider now the sequence consisting of ordered pairs

$$(6) \quad (s(d_0), \text{Min}(\omega_0)), (s(d_1), \text{Min}(\omega_1)), \dots,$$

where  $s(0) = 0$  and  $s(n) = 1$ , for  $n \geq 1$ . Because the component sequences of (6) are ultimately periodic, so is the whole sequence (6).

From this point on the proof is a straightforward modification of the proof of Lemma 5.4. in [4] (it uses only the periodicity of (5)). We omit the details.

**Remark 1.** The proof of Lemma 3 in [4] is constructive. However, our analogous proof for growing DOL-systems is not constructive, because we need the Theorem of [2].

**Remark 2.** In [4] M. Nielsen solves the growth-set equivalence problem for PDOL-systems by changing the considered PDOL-sequences effectively to strictly increasing PDOL-sequences (with the same growth-set). Lemma 2 shows that we cannot solve the growth-set equivalence problem for DOL-systems by this method. In fact, it is not known if this problem is decidable at all.

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Department of Mathematics,  
University of Turku, Finland.



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