

OPTIMAL TRANSMISSION OF GAUSSIAN SIGNALS INVOLVING COST OF FEEDBACK

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Methods to analyse the problem of signal transmission through a noiseless feedback channel using the results of nonlinear filtering theory are given in [2]. In [1] we constructed an optimal transmission model for Gaussian signals taking into account the cost of transmission. The aim of this paper is to construct an optimal transmission scheme for a Gaussian signal involving the cost of feedback.

Suppose the transmitted message is a Gaussian random variable θ with $\mathbf{E}\theta = m$ and $\mathbf{E}(\theta - m)^2 = \gamma > 0$, and the transmission of θ is carried out according to the following scheme

$$(1) \quad d\xi_t = \left\{ \beta_t [A_0(t, \xi) + A_1(t, \xi)\theta] + (1 - \beta_t) [B_0(t) + B_1(t)\theta] \right\} dt + dW_t, \\ \xi_0 = 0, \quad t \in [0, T],$$

where $\beta = (\beta_t)_{t \leq T}$ is a $(\mathcal{F}_t^\xi)_{t \leq T}$ -adapted stochastic process taking two values 0 and 1. Under $\beta_t = 1$ the transmission is with feedback and under $\beta_t = 0$ it is without feedback. Let

$$\delta(t) = \inf_{A_0, A_1, B_0, B_1, \beta, \hat{\theta}} \mathbf{E} \left[(\theta - \hat{\theta}_t)^2 + c \int_0^t \beta_s ds \right], \quad c > 0,$$

where c is the cost of the feedback.

Let the following moment conditions

$$(2) \quad \mathbf{E} \left\{ [A_0(t, \xi) + A_1(t, \xi)\theta]^2 \mid \mathcal{F}_t^\xi \right\} \leq P, \quad \mathbf{E} [B_0(t) + B_1(t)\theta]^2 \leq P$$

be satisfied, $P > 0$.

The functionals A_0 , A_1 , B_0 , B_1 , β are supposed to be such that equation (1) has a unique strong solution.

Denote

$$m_t = \mathbf{E}(\theta \mid \mathcal{F}_t^\xi), \quad \gamma_t = \mathbf{E}[(\theta - m_t)^2 \mid \mathcal{F}_t^\xi].$$

The equations for m_t and γ_t are of the following form

$$(3) \quad \begin{aligned} dm_t &= \gamma_t [\beta_t A_1(t, \xi) + (1 - \beta_t) B_1(t)] \cdot \\ &\cdot \left\{ d\xi_t - [\beta_t A_0(t, \xi) + (1 - \beta_t) B_0(t) + (\beta_t A_1(t, \xi) \right. \\ &\left. + (1 - \beta_t) B_1(t)) m_t] dt \right\}, \quad m_0 = m, \end{aligned}$$

$$(4) \quad d\gamma_t = -\gamma_t^2 [\beta_t A_1(t, \xi) + (1 - \beta_t) B_1(t)]^2 dt, \quad \gamma_0 = \gamma.$$

Theorem. Suppose the transmission of a Gaussian variable θ is carried out according to scheme (1).

If $c \geq \frac{1}{4}\gamma P$, then the optimal strategy is $\beta_s^* = 0$, $s \in [0, t]$, i.e., transmission without feedback. The optimal coding is $B_1^*(s) = \sqrt{P/\gamma}$, $B_0^*(s) = -m\sqrt{P/\gamma}$, and optimal decoding m_t^* , for given β^* , B_0^* , B_1^* , is determined through (2) and (3). In that case

$$\delta(t) = \frac{\gamma}{1 + Pt}.$$

Exactly the same statement is true if $c < \frac{1}{4}\gamma P$ and $t \leq a$ or $t \geq b$ with

$$a = \frac{\gamma P - 2c - \sqrt{\gamma^2 P^2 - 4c\gamma P}}{2cP}, \quad b = \frac{\gamma P - 2c + \sqrt{\gamma^2 P^2 - 4c\gamma P}}{2cP}.$$

If $c < \frac{1}{4}\gamma P$ and $t \in (a, b)$, then the optimal strategy has the form $\beta_s^* = I(t - x_0 \leq s \leq t)$, where x_0 is the unique solution of the equation

$$ce^{Px} [1 + P(t - x)]^2 = \gamma P^2 (t - x)$$

in the interval $(0, t)$.

The optimal coding rules are determined as follows:

$$B_1^*(s) = \sqrt{\frac{P}{\gamma}}, \quad B_0^*(s) = -m\sqrt{\frac{P}{\gamma}},$$

$$A_1^*(s, \xi) = \sqrt{\frac{P}{\gamma}} \exp \left\{ \frac{P}{2} \int_0^s \beta_u^* du \right\} \left[1 + \gamma \int_0^s (1 - \beta_u^*) (B_1^*(u))^2 \exp \left\{ -P \int_0^u \beta_r^* dr \right\} du \right],$$

$$A_0^*(t, \xi) = -A_1^*(t, \xi) m_t^*,$$

where m_t^* and γ_t^* are determined through equations (3) and (4). In this case

$$\delta(t) = \frac{\gamma e^{-Px_0}}{1 + P(t - x_0)} + cx_0.$$

To prove this theorem we need the following lemma.

Lemma. Let $\beta = (\beta_t)_{t \geq 0}$ be a real function with values in $[0, 1]$ and let P be a positive constant. Then the following inequality is valid

$$\int_0^t e^{-P \int_0^s \beta_u du} ds \leq t - \int_0^t \beta_s ds + \frac{1}{P} \left(1 - e^{-P \int_0^t \beta_s ds} \right).$$

If β has the form $\beta_s = I(u \leq s \leq t)$, $u \in [0, t]$, then the inequality reduces into an equality.

Proof. Let $\int_0^t \beta_s ds = v$ and $\beta_s^* = I(t - v \leq s \leq t)$. Then $\int_0^t \beta_s^* ds = v$. Let us first show that

$$\int_0^s \beta_u^* du \leq \int_0^s \beta_u du$$

for any $s \in [0, 1]$. This inequality is evident for $s < t - v$. Assume that for some $t_1 \geq t - v$ an inverse inequality occurs. Then

$$v = \int_0^t \beta_u^* du = \int_0^{t_1} \beta_u^* du + \int_{t_1}^t \beta_u^* du > \int_0^{t_1} \beta_u du + \int_{t_1}^t \beta_u du = \int_0^t \beta_u du,$$

which is not true. Thus,

$$\int_0^t e^{-P \int_0^s \beta_u du} ds \leq \int_0^t e^{-P \int_0^s \beta_u^* du} ds.$$

One can easily see that

$$\int_0^t e^{-P \int_0^s \beta_u^* du} ds = t - v + \frac{1}{P} (1 - e^{-Pv}),$$

and since $v = \int_0^t \beta_s ds$, we obtain the required inequality.

The validity of the second assertion of the lemma follows by a direct verification.

Proof of the theorem. For fixed coding functionals A_0, A_1, B_0, B_1 and strategy β the optimal decoding is m_t . Hence,

$$\delta(t) = \inf_{A_0, A_1, B_0, B_1, \beta} \mathbf{E} \left[\gamma_t + c \int_0^t \beta_s ds \right].$$

We construct first the optimal coding functionals for a fixed strategy β . Rewrite the moment condition in the following form

$$(5) \quad [A_0(t, \xi) + A_1(t, \xi)m_t]^2 + \gamma_t A_1^2(t, \xi) \leq P, \quad [B_0(t) + B_1(t)m]^2 + \gamma B_1^2(t) \leq P.$$

The equation for γ_t can be written in the form

$$\gamma_t = \gamma \exp \left\{ - \int_0^t \beta_s \gamma_s A_1^2(s, \xi) ds - \int_0^t (1 - \beta_s) \gamma_s B_1^2(s) ds \right\}.$$

From (5) it follows that $\gamma_t A_1^2(t, \xi) \leq P$. Therefore,

$$(6) \quad \gamma_t \geq \gamma \exp \left\{ -P \int_0^t \beta_s ds - \int_0^t (1 - \beta_s) B_1^2(s) ds \right\}.$$

If we choose $A_1^*(t, \xi) = \sqrt{P/\gamma_t^*}$, where γ_t^* is a solution of (4) for a given functional A_1^* and a fixed B_1 , then the inequality reduces into an equality. Hence A_1^* is the optimal coding. The moment condition is satisfied for $A_0^*(t, \xi) = -A_1^*(t, \xi) m_t^*$, where m_t^* is a solution of (3) for given A_0^* , A_1^* , γ_t^* . If we solve the equation (4) for $A_1^*(t, \xi) = \sqrt{P/\gamma_t^*}$, we obtain

$$\gamma_t^* = \frac{\gamma e^{-P \int_0^t \beta_s ds}}{1 + \gamma \int_0^t (1 - \beta_s) B_1^2(s) e^{-P \int_0^s \beta_u du} ds},$$

and therefore the optimal coding has the form

$$A_1^*(t, \xi) = \sqrt{\frac{P}{\gamma}} e^{P/2 \int_0^t \beta_s ds} \left[1 + \gamma \int_0^t (1 - \beta_s) B_1^2(s) e^{-P \int_0^s \beta_u du} ds \right].$$

Thus

$$\delta(t) = \inf_{B_0, B_1, \beta} \mathbf{E} \left[\frac{\gamma e^{-P \int_0^t \beta_s ds}}{1 + \gamma \int_0^t (1 - \beta_s) B_1^2(s) e^{-P \int_0^s \beta_u du} ds} + c \int_0^t \beta_s ds \right].$$

It follows from (5) that $B_1^2(t) \leq P/\gamma$. If we choose $B_1^*(t) = \sqrt{P/\gamma}$ and $B_0^*(t) = -m\sqrt{P/\gamma}$, then the condition (5) is fulfilled and

$$\delta(t) = \inf_{\beta} \mathbf{E} \left[\frac{\gamma e^{-P \int_0^t \beta_s ds}}{1 + P \int_0^t (1 - \beta_s) e^{-P \int_0^s \beta_u du} ds} + c \int_0^t \beta_s ds \right].$$

Thus, for a fixed strategy β , our construction gives the optimal coding functionals and the optimal decoding.

Now we construct the optimal strategy. It can be easily shown that

$$\delta(t) = \inf_{\beta} \mathbf{E} \left[\frac{\gamma}{1 + P e^{P \int_0^t \beta_s ds} \int_0^t e^{-P \int_0^s \beta_u du} ds} + c \int_0^t \beta_s ds \right].$$

We denote the expression in the brackets by $\delta_\beta(t)$. Using the inequality from the previous lemma we obtain

$$\delta_\beta(t) \geq \frac{\gamma e^{-P \int_0^t \beta_s ds}}{1 + P(t - \int_0^t \beta_s ds)} + c \int_0^t \beta_s ds.$$

Consider the function

$$f(x) = \frac{\gamma e^{-Px}}{1 + P(t - x)} + cx, \quad x \in [0, t].$$

The properties of the derivatives of this function easily show that this function increases for all $x \in [0, 1]$, if $c \geq \frac{1}{4}\gamma P$ and therefore $f(x) \geq f(0) = \gamma/(1 + Pt)$. If $c < \gamma P/4$ and $t \notin (a, b)$, where the numbers a and b have been defined in the enunciation of the theorem, then again $f(x) \geq \gamma/(1 + Pt)$, $x \in [0, t]$. If $c < \frac{1}{4}\gamma P$ and $t \in (a, b)$, then $f'(x) < 0$ for $x \in [0, x_0)$ and $f'(x) > 0$ for $x \in (x_0, t]$, i.e., $f(x) \geq f(x_0)$, $x \in [0, t]$, where x_0 is the unique solution of the equation

$$ce^{Px} [1 + P(t - x)]^2 = \gamma P^2(t - x).$$

Using the properties of the function f we can obtain the following statements:
If $c \geq \frac{1}{4}\gamma P$, then for any strategy β we have

$$\delta_\beta(t) \geq \frac{\gamma}{1 + Pt}, \quad (\text{note that } 0 \leq \int_0^t \beta_s ds \leq t),$$

and for $\beta_s^* = 0$, $s \in [0, t]$,

$$\delta_{\beta^*}(t) = \frac{\gamma}{1 + Pt},$$

i.e., β^* is the optimal strategy.

If $c < \frac{1}{4}\gamma P$ and $t \leq a$ or $t \geq b$, then the same results hold.

If $c < \frac{1}{4}\gamma P$ and $t \in (a, b)$, then

$$\delta_\beta(t) \geq \frac{\gamma e^{Px_0}}{1 + P(t - x_0)} + cx_0.$$

Let $\beta_s^* = I(t - x_0 \leq s \leq t)$. Then, according to the second statement in the previous lemma, we obtain

$$\begin{aligned} \delta_{\beta^*}(t) &= \frac{\gamma e^{-P \int_0^t \beta_s^* ds}}{1 + P(t - \int_0^t \beta_s^* ds)} + c \int_0^t \beta_s^* ds \\ &= \frac{\gamma e^{-Px_0}}{1 + P(t - x_0)} + cx_0, \end{aligned}$$

which completes the proof of the theorem.

References

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